International Journal of Research
Available at https://edupediapublications.org/journals
e-ISSN: 2348-6848
p-ISSN: 2348-795X
Volume 05 Issue-01
January 2018

Prevention of Sipdas Based Attacks in Cloud Computing

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Abstract--*Cloud Computing allows customers to access cloud resources and services. On-demand, self-service and pay-by-use business model are adapted for the cloud resource sharing process. Service level agreements (SLA) regulate the cost for the services that are provided for the customers. Cloud data centers are employed to share data values to the users. Denial-of-Service (DoS) attack is an attempt by attacker to prevent legitimate users from using resources. Distributed Denial of Service (DDoS) Attacks are generated in a "many to one" dimension. In DDoS attack model Large number of compromised host are gathered to send useless service requests, packets at the same time .DoS and DDoS attacks initiates the service degradation, availability and cost problems under cloud service providers.*

Brute-force attacks are raised against through specific periodic, pulsing and low-rate traffic patterns. Rate-controlling, time-window, worst-case threshold and pattern-matching are adapted to discriminate the legitimate and attacker activities. Stealthy attack patterns are raised against applications running in the cloud. Slowly-Increasing- Polymorphic DDoS Attack Strategy (SIPDAS) can be applied to initiate application vulnerabilities. SIPDAS degrades the service provided by the target application server running in the cloud. Polymorphic attacks changes the message sequence at every successive infection to avoid signature detection process. Slowly-increasing polymorphic behavior induces enough overloads on the target system. XML-based DoS (XDoS) attacks to the web-based systems are applied as the testing environment for the attack detection process.

1. INTRODUCTION

Cloud providers offer services to rent computation and storage capacity, in a way as transparent as possible, giving the impression of 'unlimited resource availability'. Such resources are not free.
Therefore, cloud providers allow customers to obtain and configure suitably the system capacity, as
well as to quickly renegotiate such capacity as their requirements change, in order that the customers can pay only for resources that they actually use. Several
cloud providers offer the 'load balancing' service for automatically distributing the incoming application service requests across multiple instances, as well as the 'auto scaling' service for enabling consumers to closely follow the demand curve for their applications. In order to minimize the customer costs, the auto scaling ensures that the number of the application instances increases seamlessly during the demand spikes and decreases automatically during the demand lulls. For example, by using Amazon EC2 cloud services, the consumers can set a condition to add new computational instances when the average CPU utilization exceeds a fixed threshold. Moreover, they can configure a cool-down period in order to allow the application workload to stabilize before the auto scaling adds or removes the instances. In the following, we will show how this feature can be maliciously exploited by a stealthy attack, which may slowly exhaust the resources provided by the cloud provider for ensuring the SLA, and enhance the costs incurred by the cloud customer.

2. RELATED WORK

We briefly outline some masquerade detection approaches. The uniqueness approach assumes that commands that have not been seen in the training data indicate a masquerader. Moreover, the probability that a masquerader has issued a command is inversely related to the number of users that use such a command. While uniqueness has a relatively poor performance, it is one of the few approaches that target false alarm rate of 1 percent. Naive Bayes One-step Markov is based upon one-step transitions from a command to the next. It builds two transition matrices for each user from, respectively, the training database and the testing one and it triggers an alarm when these matrices noticeably differ. The false

alarm rate of this method is not satisfactory.

The Hybrid Multi-Step Markov method is based on Markov chains. When a Markov model cannot be adopted because too many commands in the testing data have not been observed in the training, a simple independence model with probabilities estimated from a contingency table of users versus commands may be more appropriate. Schonlau et al. toggled between a Markov model and the simple independence one. This approach achieves the best performance among the considered methods. The main idea underlying the compression approach is that new and old data from the same user should compress at about the same ratio. Instead, data from a masquerading user will compress at a different ratio. Among the proposed methods, this results in the worst performance. Incremental Probabilistic Action Modeling (IPAM) is based upon one-step command transition. It estimates the probability of each transition from the training data set and uses it to predict the sequence of user commands. Too many false predictions signal a masquerader. This method is in the lowest-performing group. Sequence-matching computes a similarity match between the user profiles and the corresponding sequence of commands. Any score lower than a threshold signals a masquerader. Its performance on the SEA data set is not very high.

Support Vector Machine (SVM) denotes a set of machine learning algorithms for binary data classification. It exploits a set of support vectors in the training data that outlines a hyper plane in feature space. SVM can potentially learn a large set of patterns but it results in high false alarm rates and a low detection rate. Furthermore, the user profile has to be updated to reduce false alarms. Szymanski and Zhang propose a recursive data mining approach that discovers frequent patterns in the sequence of user commands, encodes them with unique symbols, and rewrites the sequence with the new coding. Then, a one-class SVM classifier detects masqueraders. This approach demands mixing user data and may not be ideal or easily implemented in real-world. It also suffers of some of the SVM shortcomings. Maxion and Townsend applied a Naive Bayes classifier widely used in text classification tasks and that classifies sequences of user-command data into either legitimate or masquerader. The method has not yet achieved the level of accuracy for practical deployment. Dash et al. introduced an episode based Naive Bayes technique that extracts meaningful episodes from a long sequence of commands.

The Naive Bayes algorithm identifies these episodes either as masquerade or normal according to the number of commands in masquerade blocks. The proposed technique significantly improves the hit ratio but it still has high false positive rates and it does not update the user profile. Alok et al. integrates a Naive Bayes approach with one based on a weighted radial basis function, WRBF, similarity. The Naive Bayes algorithm includes information on the probabilities of commands by one user over the other users. Instead, the WRBF similarity takes into account the similarity measure based on the frequency of commands, f, and the weight associated with the frequency vectors. Here, f is a similarity score between an input frequency vector and a frequency vector from the training data set. The experiments confirm that WRBF-NB significantly improves the hit ratio but, as the previous approach, it suffers from the high false positive rates. Furthermore, it increases the overall overhead by computing both the Naive Bayes and the WRBF and integrating their results. Lastly, it does not update the user profile and neglects the low level representation of user commands. Dash et al. introduced an adaptive Naive Bayes approach based on the premise that both the commands of a legitimate user and those of an attacker may differ from the trained signature but the deviation of the legitimate user is momentary, whereas the attacker one persists longer. The improvement in the performance of detection has been empirically verified using several data sets. The false positive rate is still high.

Malek and Salvatore have modeled user OS commands as bag-of-words without timing information. They used a one-class support vector machine to achieve a better performance than threshold based comparison with a distance metric. The ability of sequence alignment algorithms to find areas of similarity can be exploited to differentiate legitimate usage from masquerade attacks. To do so, a signature of the normal user behavior should be created by collecting sequences of audit data. Then, this signature is aligned with audit data from monitored sessions to find areas of similarity. Areas that do not align properly are assumed to be anomalous, and several anomalous areas are a strong indicator of masquerade attacks. Among sequence alignment algorithms such as global, local and semi-global alignments, the most efficient one is semi-global alignment. Adesina et al. modified the scoring function of the semi-global alignment algorithm to improve the detection efficiency.

They used a systematically generated ASCII coded sequence audit data from Windows and UNIX systems as simulations for the intrusion data set. A real time evaluation using one of the current data sets is missing.

Coull et al. modified the Smith-Waterman alignment algorithm a semi-global alignment algorithm that is described with some evolutions and enhancements. Different techniques in terms of the Receiver Operator Characteristic (ROC) curves and the Maxion-Townsend cost function and it shows that SGA achieves the best performances. The Maxion-Townsend cost function rates a masquerade detection algorithm and defines the detection cost.

3. HANDLING DENIAL OF SERVICE ATTACKS IN CLOUD

Cloud Computing is an emerging paradigm that allows customers to obtain cloud resources and services according to an on-demand, self-service, and pay-by use business model. Service level agreements (SLA) regulate the costs that the cloud customers have to pay for the provided quality of service (QoS). A side effect of such a model is that, it is prone to Denial of Service (DoS) and Distributed DoS (DDoS), which aim at reducing the service availability and performance by exhausting the resources of the service's host system. Such attacks have special effects in the cloud due to the adopted pay-by-use business model. Specifically, in cloud computing also partial service degradation due to an attack has direct effect on the service costs, and not only on the performance and availability perceived by the customer. The delay of the cloud service provider to diagnose the causes of the service degradation can be considered as security vulnerability. It can be exploited by attackers that aim
at exhausting the cloud resources and seriously degrading the QoS, as happened to the BitBucket Cloud, which went down for 19h. Therefore, the cloud management system has to implement specific countermeasures in order to avoid paying credits in case of accidental or deliberate intrusion that cause violations of QoS guarantees.

Over the past decade, many efforts have been devoted to the detection of DDoS attacks in distributed systems. Security prevention mechanisms usually use approaches based on rate-controlling, time-window, worst-case threshold, and pattern-matching methods to discriminate between the nominal system operation and malicious behaviors. On the other hand, the attackers are aware of the presence of such protection mechanisms. They attempt to perform their activities in a "stealthy" fashion in order to elude the security mechanisms, by orchestrating and timing attack patterns that leverage specific weaknesses of target systems. They are carried out by directing flows of legitimate service requests against a specific system at such a low-rate that would evade the DDoS detection mechanisms, and prolong the attack latency, i.e., the amount of time that the ongoing attack to the system has been undetected.

This paper presents a sophisticated strategy to orchestrate stealthy attack patterns against applications running in the cloud. Instead of aiming at making the service unavailable, the
proposed strategy aims at exploiting the cloud flexibility, forcing the application to consume more resources than needed, affecting the cloud customer more on financial aspects than on the service availability. The attack pattern is orchestrated in order to evade, greatly delay the techniques proposed in the literature to detect low-rate attacks. It does not exhibit a periodic waveform typical of low-rate exhausting attacks. In contrast with them, it is an iterative and incremental process. In particular, the attack potency is slowly enhanced by a patient attacker, in order to inflict significant financial losses, even if the attack pattern is performed in accordance to the maximum job size and arrival rate of the service requests allowed in the system. Using a simplified model empirically designed, we derive an expression for gradually increasing the potency of the attack, as a function of the reached service degradation. We show that the features offered by the
cloud provider, to ensure the SLA negotiated with the customer can be maliciously exploited by the proposed. Stealthy attack, which slowly exhausts the resources provided by the cloud provider and increases the costs incurred by the customer.

The proposed attack strategy, namely Slowly-Increasing-Polymorphic DDoS Attack Strategy (SIPDAS) can be applied to several kind of attacks, that leverage known application vulnerabilities, in order to degrade the service provided by the target application server running in the cloud. The term polymorphic is inspired to polymorphic attacks which change message sequence at every successive infection in order to evade signature detection mechanisms. Even if the victim detects the SIPDAS attack, the attack strategy can be reinitiate by using a

different application vulnerability, or a different timing.

In order to validate the stealthy characteristics of the proposed SIPDAS attack, we explore potential solutions proposed in the literature to detect sophisticated low-rate DDoS attacks. We show that the proposed slowly-increasing polymorphic behavior induces enough overload on the target system and evades, or however, delays greatly the detection methods. In order to explore the attack impact against an application deployed in a cloud environment, this paper focuses on one of the most serious threats to cloud computing, which comes from XML-based DoS (XDoS) attacks to the web-based systems. The experimental testbed is based on the mOSAIC framework, which offers both a 'Software Platform' that enables the execution of applications developed using the mOSAIC API, and a 'Cloud Agency', that acts as a provisioning system, brokering resources from a federation of cloud providers.

4. PROBLEM STATEMENT

Brute-force attacks are raised against through specific periodic, pulsing and low-rate traffic patterns. Rate-controlling, time-window, worst-case threshold and pattern-matching are adapted to discriminate the legitimate and attacker activities. Stealthy attack patterns are raised against applications running in the cloud. Slowly-Increasing-
Polymorphic DDoS Attack Strategy (SIPDAS) can be applied to initiate application vulnerabilities.
SIPDAS degrades the service provided by the target application server running in the cloud. Polymorphic attacks changes the message sequence at every successive infection to avoid signature detection process. Slowly-increasing polymorphic behavior induces enough overloads on the target system. XML-based DoS (XDoS) attacks to the web-based systems are applied as the testing environment for the attack detection process. The following drawbacks are identified from the existing system.

- SIPDAS based attack detection is not supported
- Polymorphic behavior identification is not adapted
- Application level vulnerability detection is low
- Service degradation and resource consumption cost analysis is not performed

5. STEALTHY DOS ATTACKS ON CLOUD SERVICES

5.1. DoS Attacks Against Cloud Applications

In this section are presented several attack examples, which can be leveraged to implement the proposed SIPDAS attack pattern against a cloud application. In particular, we consider DDoS attacks that exploit application vulnerabilities, including: the Oversize Payload attack that exploits the high memory consumption of XML processing; the Oversized Cryptography that exploits the flexible usability of the security elements defined by the WS-Security specification the Resource Exhaustion attacks use flows of messages that are correct regarding their message structure, but that are not properly correlated to any existing process instance on the target server and attacks that exploit the worst-case performance of the system, for example by achieving the worst case complexity of Hash table data structure, or by using complex queries that force to spend much CPU time or disk access time.

In this paper, we use a Coercive Parsing attack as a case study, which represents one of the most serious threat for the cloud applications. It exploits the XML verbosity and the complex parsing process. In particular, the Deeply-Nested XML is a resource exhaustion attack, which exploits the XML message format by inserting a large number of nested XML tags in the message body. The goal is to force the XML parser within the application server, to exhaust the computational resources by processing a large number of deeply-nested XML tags.

5.2. Stealthy Attack Objectives

The system is aimed to defining the objectives that a sophisticated attacker would like to achieve, and the requirements the attack pattern has to satisfy to be stealth. Recall that, the purpose of the attack against cloud applications is not to necessarily deny the service, but rather to inflict significant degradation in some aspect of the service, namely attack profit PA, in order to maximize the cloud resource consumption CA to process malicious requests. In order to elude the attack detection, different attacks that use low-rate traffic have been presented in the literature. Therefore, several works have proposed techniques to detect low-rate DDoS attacks, which monitor anomalies in the fluctuation of the incoming traffic through either a time or frequency-domain analysis. They assume that, the main anomaly can be incurred during a low-rate attack is that, the incoming service requests fluctuate in a more extreme

manner during an attack. The abnormal fluctuation is a combined result of two different kinds of behaviors: (i) a periodic and impulse trend in the attack pattern, and (ii) the fast decline in the incoming traffic volume. Therefore, in order to perform the attack in stealthy fashion with respect to the proposed detection techniques, an attacker has to inject low-rate message flows $\varphi_{Aj} = [\varphi_{j,1}, \ldots, \varphi_{j,m}]$, that satisfy the following optimization problem:

5.3. Attack Approach

In order to implement SIPDAS-based attacks, the following components are involved:

- a Master that coordinates the attack;
- π Agents that perform the attack; and
- a Meter that evaluates the attack effects.

The approach implemented by each Agent to perform a stealthy service degradation in the cloud computing. It has been specialized for an X-DoS attack. Specifically, the attack is performed by injecting polymorphic bursts of length T with an increasing intensity until the attack is either successful or detected. Each burst is formatted in such a way as to inflict a certain average level of load C_R. In particular, we assume that C_R is proportional to the attack intensity of the flow Φ_{Aj} during the period T. Therefore, denote I_0 as the initial intensity of the attack, and assuming $\Delta C_R = \Delta I$ as the increment of the attack intensity. For each attack period, fixed the maximum number of nested tags (tagThreshold), the routine pick RandomTags(. . .) randomly returns the number of nested tags nT for each message. Based on nT , the routine compute Inter arrival Time uses a specific algorithm to compute the inter-arrival time for injecting the next message. At the end of the period T, if the condition 'attack Successful' is false, the attack intensity is increased. If the condition 'attack Successful' is true, the attack intensity is maintained constant until either the attack is detected or the auto-scaling mechanism enabled in the cloud adds new cloud resources. The attack is performed until it is either detected, or the average message rate of the next burst to be injected is greater than dT. In this last case, the Agent notifies to the Master that the maximum average message rate is reached and continues to inject messages formatted according to the last level of load CR reached.

6. CONCLUSION

In this paper, we have a tendency to propose a technique to implement sneaky attack patterns, that exhibit a slowly-increasing polymorphic behavior that may evade, or however, greatly delay the techniques planned within the literature to find low-rate attacks. Exploiting a vulnerability of the target application, a patient and intelligent assailant will orchestrate subtle flows of messages, indistinguishable from legitimate service requests above all, the planned attack pattern, rather than aiming at creating the service inaccessible, it aims at exploiting the cloud flexibility, forcing the services to rescale and consume additional resources than required, affecting the cloud client additional on money aspects than on the service availableness. within the future work, we have a tendency to aim at extending the approach to a bigger set of application level vulnerabilities, likewise as process a classy methodology able to find SIPDAS based attacks within the cloud computing setting.

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