
An Efficient Dynamic Job Ordering and Slot Configuration for Minimizing the Makespan of MapReduce Jobs

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Abstract: MapReduce is a well-known parallel computing paradigm for massive data processing in clusters and data facilities. It is observed that different job execution orders and MapReduce slot configurations for a MapReduce workload have the enormously extraordinary efficiency related to the makespan, complete completion time, process utilization and other performance metrics. This paper proposes two classes of algorithms to diminish the lifespan and the total completion time for an offline MapReduce workload. Our first class of algorithms focuses on the job ordering optimization for a MapReduce workload below a given map/reduce slot configuration. In distinction, our 2nd type of algorithms considers the scenario that we will participate in optimization for map/reduce down slot configuration for a MapReduce workload.

Keywords- MapReduce, Hadoop, Flow-shops, Scheduling algorithm, Job ordering

I. INTRODUCTION

A MapReduce job consists of a set of map and reduce tasks, where reduce tasks are performed after the map tasks. Hadoop [2], an open source implementation of MapReduce, has been deployed in large clusters containing thousands of machines by companies such as Amazon and Facebook. Make span and total completion time are two key performance metrics. Generally, make span is defined as the time period since the start of the first job until the completion of the last job for a set of jobs. It considers the computation time of jobs and is often used to measure the performance and utilization efficiency of a system. In contrast, total completion time is referred to as the sum of completed time periods for all jobs since the start of the first job.

It is a generalized make span with queuing time (i.e., waiting time) included. We can use it to measure the satisfaction to the system from a single job's perspective through dividing the total completion time by the number of jobs (i.e., average completion time).

In those cluster and data center environments, MapReduce and Hadoop are used to support batch processing for jobs submitted from multiple users (i.e., MapReduce workloads). Despite many research efforts devoted to improve the performance of a single MapReduce job. There are two key performance metrics i.e. Makespan and total completion time (TCT) and we aim to optimize these metrics. Generally, make span is defined as the time period since the start of the first job until the completion of the last job for a set of jobs. It considers the computation time of jobs and is often used to measure the performance and utilization efficiency of a system. In contrast, total completion time is referred to as the sum of completed time periods for all jobs since the start of the first job. It is a generalized makespan with queuing time (i.e., waiting time) included. We can use it to measure the satisfaction to the system from a single job's perspective through dividing the total completion time by the number of jobs (i.e., average completion time). Therefore, in this paper, we aim to optimize these two metrics the number of jobs (i.e., average completion time). Therefore, in this paper, we aim to optimize these two metrics.

Objectives:

- To improve the performance for MapReduce workloads with job ordering and slot configuration optimization approaches.

- Propose slot configuration algorithms for make span and total completion time.
- Perform extensive experiments to validate the effectiveness of proposed algorithms and theoretical results

II. RELATED WORK

The important problems are that we define and more issue in scheduling technique of MapReduce, the scheduling is one of the most critical criteria of MapReduce. There are many algorithm can address these problem with different techniques and methods. Some of them focus on dynamic slot allocation and straggler problem for speculative execution. Also many of them have been design deadline constrain to minimizing the total job completion time. In this section we describe some of these algorithm techniques.

Wolf et al. [2] implemented flexible scheduling allocationscheme with Hadoop fair scheduler. A primary concern is to optimize scheduling theory metrics, response time, makespan, stretch, and Service Level Agreement. They proposed penalty function for measurement of job completion time, epoch scheduling for partitioning time, moldable scheduling for job parallelization, and malleable scheduling for different interval parallelization.

Dean et al. 2008 [1] have discussed MapReduce programming model. The MapReduce model performs operations using the map and reduces functions. Map function gets input from user documents. It generates intermediate key/value for reducing function. It further processes intermediate key/value pairs and provide output key/value pairs. At an entry level, MapReduce programming model provided the best data processing results. Currently, it needs to process the large volume of data. So it provides some consequences while processing and generating data sets. It takes much execution time for task initialization, task coordination, and task scheduling. Parallel data processing may lead to inefficient task execution and low resource utilization.

Verma et al. [3] proposed two algorithms for makespan optimization. First is a greedy algorithm job ordering method based on Johnson's Rule. Another is a heuristic algorithm called BalancedPool. They have introduced a simple abstraction where each MapReduce job is represented as a pair of map and reduce stage duration. The Johnson algorithm was designed for building an optimal job schedule. This framework evaluates the performance benefits of the constructed schedule through an extensive set of simulation over a variety of realistic workloads. It measures how many numbers of slots required for scheduling the slots dynamically with a particular job deadline.

Tang et al. [4] have proposed three techniques to improve MapReduce performance. First technique is Dynamic Hadoop Slot Allocation. They categorized utilized slot into the busy slot and idle slot respectively. The primary concern is to increase the number of the busy slots and decrease number of idle slots. DHSA observes idle map and reduce slots. Dynamic Hadoop Slot Allocation allocate the task only to the unallocated map slots and due to Speculative Execution Performance Balancing provides performance upgrade for a batch of jobs. It gives the highest priority to failed tasks and next level priority to pending tasks. Due to slot pre scheduling it improves the performance of slot utilization.

Tang, Lee and He [5] have proposed DynamicMR: A Dynamic Slot Allocation Optimization Framework for improving the performance for a single job but at the expense of the cluster efficiency. They proposed Hazardous Execution Performance Balancing technique for balancing the performance tradeoff between a single job and a batch of jobs. Slot PreScheduling is the new technique and that can improve the data locality but with no impact on fairness. Finally, integrating these two techniques, new technique is implemented called DynamicMR that can improve the performance of MapReduce workloads. Tang, Lee and He [6] have proposed MROrder: Flexible Job Ordering technique which optimizes the job order for online MapReduce workloads. MROrder is designed to be flexible for

different optimization metrics, e.g., makespan and total completion time

The computing Resources are considering into map and reduce slots, which are primary computing units and statically configure by administrator in advance. A MapReduce job execution has two main features: first, the slot allocation constrains that map slots allocate to map slots and reduce slots can be allocated to reduce task. There is immense different performance and system utilization for a MapReduce over differ slot configurations. So if we use each and every slot in both slot according to needs of node which affects the system utilization and performance. For that use a Dynamic Hadoop Slot Allocation (DHSA) [7] technique to increase the slot utilization for MapReduce. Straggler Problem occur because of unavoidable run-time contention for processor, memory, network bandwidth and also other resources causing great affect on delay of the whole job. For that use a Speculative Execution Performance Balancing (SEPB) [7] to balance the utilization for single job as well as multiple jobs.

Data Locality increase slot utilization efficiency and performance achieve great output for MapReduce workloads. After all, there is often struck between fairness and data optimization in a shared cluster among many users. For that use a Slot PreScheduling technique to achieve great significant data locality at the expense of the load balance nodes. User having specific job deadline so deadline are most important requirement which can improve the performance. MapReduce Task Scheduling algorithm for Deadline Constrains (MTSD) [8] has main goal on user's deadline constraints problem.

III. PROPOSED WORK

(Algorithm1) Slot allocation and Slot pre-scheduling process:

In this module, we are going to perform two processes. Slot allocation Slot pre-scheduling process. In this slot allocation process, we are going to allocate the slot based on dynamic Hadoop slot allocation optimization mechanism. In the slot pre-scheduling process we are going to improve the data

locality. Slot Pre-Scheduling technique that can improve the data locality while having no negative impact on the fairness of Map-Reduce jobs. Some idle slots which cannot be allocated due to the load balancing constraint during runtime, we can pre-allocate those slots of the node to jobs to maximize the data locality.

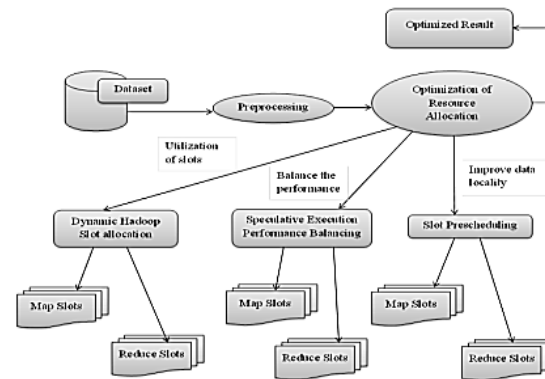


Fig.1 Overview of Dynamic Job Ordering

(Algorithm2) Speculative Execution Performance Balancing:

When a node has an idle map slot, we should choose pending map tasks first before looking for speculative map tasks for a batch of jobs. Hadoop Slot is executed for determining the path for performing the MapReduce job. After this, the Speculative based process starts to execute the determined optimized Multi-execution path. Executing individual MapReduce jobs in each datacenter on corresponding inputs and then aggregating results is defined as a MULTI execution path. This path used to execute the jobs effectively.

IV. CONCLUSION

Dynamic slot configuration is one of the significant factors while processing a large data set with MapReduce model. It optimizes the enactment of MapReduce framework. Each job can be scheduled using any one of the scheduling policies by the job tracker. The task managers which are present in the task tracker allocate slots to jobs. As of the inspected paper, it is concluded to prefer a dynamic slot allocation strategy that includes active jobs workload

estimation, optimal slot assignment, and scheduling policy.

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BIODATA

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