

# A Duplicates New Approach for Growth the Performance of Locating

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## ABSTRACT:

*With the ever increasing volume of data, data quality problems abound. Multiple, yet different representations of the same real-world objects in data, duplicates, are one of the most intriguing data quality problems. The effects of such duplicates are detrimental. For instance, bank customers can obtain duplicate identities, inventory levels are monitored incorrectly, catalogs are mailed multiple times to the same household, etc. Automatically detecting duplicates is difficult. Duplicate detection is the process for identifying multiple representations of same real world entities. Nowadays, duplicate detection methods need to*

*process ever larger datasets in ever shorter time: maintaining the quality of a dataset becomes increasingly difficult. Genetic algorithm is proposed that significantly increase the efficiency of finding duplicates if the execution time is limited. This efficiently detects the text document duplication which has same content with distinct file name or different content with same file name*

## 1 INTRODUCTION

With the explosive growth of digital data, deduplication techniques are widely employed to backup data and minimize network and storage overhead by detecting and eliminating redundancy among data. Instead of keeping multiple data copies with the same content, deduplication eliminates redundant data by keeping only one physical

copy and referring other redundant data to that copy. Deduplication has received much attention from both academia and industry because it can greatly improve storage utilization and save storage space, especially for the applications with high deduplication ratio such as archival storage systems. A number of deduplication systems have been proposed based on various deduplication strategies such as client-side or server-side deduplications, file-level or block-level deduplications. A brief review is given in Section 6. Especially, with the advent of cloud storage, data deduplication techniques become more attractive and critical for the management of ever-increasing volumes of data in cloud storage services which motivates enterprises and organizations to outsource data storage to third-party cloud providers, as evidenced by many real-life case studies [1]. According to the analysis report of IDC, the volume of data in the world is expected to reach 40 trillion gigabytes in 2020 [2]. Today's commercial cloud storage services, such as Dropbox, Google Drive and Mozy, have been applying deduplication to save the network bandwidth and the storage cost with client-side deduplication. There are two types of

deduplication in terms of the size: (i) *file-level deduplication*, which discovers redundancies between different files and removes these redundancies to reduce capacity demands, and (ii) *blocklevel deduplication*, which discovers and removes redundancies between data blocks. The file can be divided into smaller fixed-size or variable-size blocks. Using fixedsize blocks simplifies the computations of block boundaries, while using variable-size blocks (e.g., based on Rabin fingerprinting [3]) provides better deduplication efficiency.

### 1.1 Our Contributions

In this paper, we show how to design secure deduplication systems with higher reliability in cloud computing. We introduce the distributed cloud storage servers into deduplication systems to provide better fault tolerance. To further protect data confidentiality, the secret sharing technique is utilized, which is also compatible with the distributed storage systems. In more details, a file is first split and encoded into fragments by using the technique of secret sharing, instead of encryption mechanisms. These shares will be distributed across multiple independent storage servers. Furthermore, to support deduplication, a

short cryptographic hash value of the content will also be computed and sent to each storage server as the fingerprint of the fragment stored at each server. Only the data owner who first uploads the data is required to compute and distribute such secret shares, while all following users who own the same data copy do not need to compute and store these shares any more. To recover data copies, users must access a minimum number of storage servers through authentication and obtain the secret shares to reconstruct the data. In other words, the secret shares of data will only be accessible by the authorized users who own the corresponding data copy.

## 2 PROBLEM FORMULATION

### 2.1 System Model

This section is devoted to the definitions of the system model and security threats. Two kinds entities will be involved in this deduplication system, including the user and the storage cloud service provider (S-CSP). Both client-side deduplication and server-side deduplication are supported in our system to save the bandwidth for data uploading and storage space for data storing.

- *User*. The user is an entity that wants to outsource data storage to the S-CSP and

access the data later. In a storage system supporting deduplication, the user only uploads unique data but does not upload any duplicate data to save the upload bandwidth. Furthermore, the fault tolerance is required by users in the system to provide higher reliability.

- *S-CSP*. The S-CSP is an entity that provides the outsourcing data storage service for the users. In the deduplication system, when users own and store the same content, the S-CSP will only store a single copy of these files and retain only unique data. A deduplication technique, on the other hand, can reduce the storage cost at the server side and save the upload bandwidth at the user side. For fault tolerance and confidentiality of data storage, we consider a quorum of S-CSPs, each being an independent entity. The user data is distributed across multiple S-CSPs. We deploy our deduplication mechanism in both file and block levels. Specifically, to upload a file, a user first performs the file-level duplicate check. If the file is a duplicate, then all its blocks must be duplicates as well, otherwise, the user further performs the blocklevel duplicate check and identifies the unique blocks to be uploaded. Each data

copy (i.e., a file or a block) is associated with a *tag* for the duplicate check. All data copies and tags will be stored in the S-CSP.

## 2.2 Threat Model and Security Goals

Two types of attackers are considered in our threat model: (i) An outside attacker, who may obtain some knowledge of the data copy of interest via public channels.

An outside attacker plays the role of a user that interacts with the S-CSP; (ii) An inside attacker, who may have some knowledge of partial data information such as the ciphertext. An insider attacker is assumed to be honest-but-curious and will follow our protocol, which could refer to the S-CSPs in our system. Their goal is to extract useful information from user data. The following security requirements, including confidentiality, integrity, and reliability are considered in our security model.

*Confidentiality.* Here, we allow collusion among the SCSPs. However, we require that the number of colluded S-CSPs is not more than a predefined threshold. To this end, we aim to achieve data confidentiality against collusion attacks. We require that the data distributed and stored among the S-CSPs remains secure when they are unpredictable (i.e., have high min-entropy), even if the

adversary controls a predefined number of S-CSPs. The goal of the adversary is to retrieve and recover the files that do not belong to them. This requirement has recently been formalized in [6] and called the privacy against chosen distribution attack. This also implies that the data is secure against the adversary who does not own the data. *Integrity.* Two kinds of integrity, including tag consistency and message authentication, are involved in the security model. Tag consistency check is run by the cloud storage server during the file uploading phase, which is used to prevent the duplicate/ciphertext replacement attack. If any adversary uploads a maliciously-generated ciphertext such that its tag is the same with another honestly-generated ciphertext, the cloud storage server can detect this dishonest behavior. Thus, the users do not need to worry about that their data are replaced and unable to be decrypted. Message authentication check is run by the users, which is used to detect if the downloaded and decrypted data are complete and uncorrupted or not. This security requirement is introduced to prevent the insider attack from the cloud storage service providers.

*Reliability.* The security requirement of reliability in deduplication means that the storage system can provide fault tolerance by using the means of redundancy. In more details, in our system, it can be tolerated even if a certain number of nodes fail. The system is required to detect and repair corrupted data and provide correct output for the users.

### **EXISTINGSYSTEM:**

Databases play an important role in today's IT-based economy. Many industries and systems depend on the accuracy of databases to carry out operations. Therefore, the quality of the information (or the lack thereof) stored in the databases can have significant cost implications to a system that relies on information to function and conduct business. Much research on duplicate detection, also known as entity resolution and by many other names focuses on pair selection algorithms that try to maximize recall on the one hand and efficiency on the other hand. Adaptive

techniques are capable of estimating the quality of comparison candidates. The algorithms use this information to choose the comparison candidates more carefully. In the last few years, the economic need for progressive algorithms also initiated some concrete studies in this domain. For instance, pay-as-you-go algorithms for information integration on large scale datasets have been presented. Other works introduced progressive data cleansing algorithms for the analysis of sensor data streams. However, these approaches cannot be applied to duplicate detection.

### **DISADVANTAGES OF EXISTING SYSTEM:**

- These adaptive techniques dynamically improve the efficiency of duplicate detection, but in contrast to our progressive techniques, they need to run for certain periods

of time and cannot maximize the efficiency for any given time slot

- Needs to process large dataset in short time
- Quality of data set becomes increasingly difficult

### **PROPOSED SYSTEM:**

In an error-free system with perfectly clean data, the construction of a comprehensive view of the data consists of linking—in relational terms, joining—two or more tables on their key fields. Unfortunately, data often lack a unique, global identifier that would permit such an operation. Furthermore, the data are neither carefully controlled for quality nor defined in a consistent way across different data sources. Thus, data quality is often compromised by many factors. In this proposed system two novel, progressive duplicate detection algorithms namely progressive sorted

neighborhood method (PSNM), which performs best on small and almost clean datasets, and progressive blocking (PB), which performs best on large and very dirty datasets. Both enhance the efficiency of duplicate detection even on very large datasets. In this project genetic programming algorithm is used to detect the duplication of text document. Text document which has same content with different name is detected and saved as duplicate. If any document has same name with different content is saved without overwrite.

### **ADVANTAGES OF PROPOSED SYSTEM:**

#### **□ Improved early quality**

Let  $t$  be an arbitrary target time at which results are needed. Then the progressive algorithm discovers more duplicate pairs at  $t$  than the corresponding traditional algorithm.

Typically,  $t$  is smaller than the overall runtime of the traditional algorithm.

**□ Same eventual quality** If both a traditional algorithm and its progressive version finish execution, without early termination at  $t$ , they produce the same results.

## CONCLUSIONS

We proposed the distributed deduplication systems to improve the reliability of data while achieving the confidentiality of the users' outsourced data without an encryption mechanism. Four constructions were proposed to support file-level and fine-grained block-level data deduplication. The security of tag consistency and integrity were achieved. We implemented our deduplication systems using the Ramp secret sharing scheme and demonstrated that it incurs small encoding/decoding overhead compared to the network transmission overhead in regular upload/download operations.

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