



Clone Detection In Wsn's For Optimization Of Energy And Memory Efficiency

Yelagandula Shailaja , .Dr. R. Naveen kumar , .V.Janaki

- 1.PG Scholar, Department of CSE ,Vaagdevi College of Engineering, Bollikunta, Warangal, Telangana Mail id :shailu.evergreen@gmail.com
2. Professor Department of CSE, Vaagedevi College of Engineering, Bollikunta, Warangal, Telangana Mail id :Naveensmitha@gmail.com
3. Professor,HOD Department of CSE, Vaagedevi College of Engineering, Bollikunta, Warangal,Telangana

ABSTRACT

The Energy-Efficient location-aware clone detection protocol in densely deployed WSNs, which canguarantee successful clone attack detection and maintain satisfactory network lifetime. Specifically, we exploit the location informationof sensors and randomly select witnesses located in a ring area to verify the legitimacy of sensors and to report detected clone attacks.The ring structure facilitates energy-efficient data forwarding along the path towards the witnesses and the sink. We theoretically provethat the proposed protocol can achieve 100 percent clone detection probability with trustful witnesses. We further extend the work bystudying the clone detection performance with untrustful witnesses and show that the clone detection probability still approaches98 percent when 10 percent of witnesses are compromised.

I INTRODUCTION

Moreover, in most existing clone detection protocols with random witnessselection scheme, the required buffer storage of sensors is usually dependent on the node density, protocol, the required bufferstorage of sensors is independent of n but a function of the hop length of the network radius h , i.e.,Extensive simulations demonstrate that our proposed protocol can achieve long network lifetime by effectively distributing the trafficload across the network.Wireless sensors have been widely deployed for a varietyof applications, ranging from environment monitoring to tele medicine and objects tracking, etc. [2]–[4]. For costeffective sensor placement, sensors are usually not tamper proofdevices and are deployed in places without monitoring and protection, which makes them prone to different attack-s . For example, a malicious user may compromise some sensors and acquire their privateinformation. Then, it can duplicate the

sensors and deploy clones in a wireless sensor network (WSN) to launch a variety of attacks [10], which is referred to as the clone attack [11]–[13]. As the duplicated sensors have the same information, e.g., code and cryptographic information, captured from legitimate sensors, they can easily participate in network operations and launch attacks. Due to the low cost for sensor duplication and deployment, clone attacks have become one of the most critical security issues in WSNs. Thus, it is essential to effectively detect clone attacks in order to ensure healthy operation of WSNs. To allow efficient clone detection, usually, a set of nodes are selected, which are called witnesses, to help certify the legitimacy of the nodes in the network. The private information of the source node, i.e., identity and the location information are shared with witnesses at the stage of witness selection. When any of the nodes in the network wants to transmit data, it first sends the request to the witnesses for legitimacy verification, and witnesses will report a detected attack if the node fails the certification. To achieve successful clone detection, witness selection and legitimacy verification.

II RELATED WORK

As one of the utmost important security issues, clone attack has attracted people's attention. There are many works [14],[15], [16] that studies clone detection protocols in the

literature, which can be classified into two different categories, i.e., centralized and distributed clone detection protocols. In centralized protocols, the sink or witnesses generally locate in the center of each region, and store the private information of sensors. When the sink or witnesses receive the private information of the source node, they can determine whether there is a clone attack by comparing the private information with its pre-stored records. Normally, centralized clone detection protocols have low overhead and running complexity. However, the security of sensors' private information may not be guaranteed, because the malicious users can eavesdrop the transmission between the sink node and sensors. Moreover, the network lifetime may be dramatically decreased since the sensor nodes close to the sink will deplete their energy sooner than other nodes. Different from centralized protocols, in distributed clone detection protocols, a set of witnesses are selected to match with every sensor [10], [11], which prevents the transmission between the sink and sensors from being eavesdropped by malicious users. There are three different types of witness selection schemes in distributed clone detection protocols: i) deterministic selection, ii) random selection, and iii) semi-random selection. The deterministic witness selection based clone



detection protocols like RED [10] choose the same set of witnesses for all sensor nodes. By using deterministic witness selection, a low communication overhead and a high clone detection probability can be achieved.

III Energy and Memory Efficient Clone Detection in Wireless Sensor Networks

In the literature, some distributed clone detection protocols have been proposed, such as Randomized Efficient and Distributed protocol (RED) [10] and Line-Select Multi-cast protocol (LSM) [11]. However, most approaches mainly focus on improving clone detection probability without considering efficiency and balance of energy consumption in WSNs. With such kind of approaches, some sensors may use up their batteries due to the unbalanced energy consumption, and dead sensors may cause network partition, which may further affect the normal operation of WSNs. Christo Ananth et al. [3] discussed about a system, In this proposal, a neural network approach is proposed for energy conservation routing in a wireless sensor network. Our designed neural network system has been successfully applied to our scheme of energy conservation. Neural network is applied to predict Most Significant Node and selecting the Group Head amongst the association of sensor nodes in the network. After having a precise prediction about

Most Significant Node, we would like to expand our approach in future to different WSN power management techniques and observe the results. In this proposal, we used arbitrary data for our experiment purpose; it is also expected to generate a real time data for the experiment in future and also by using adhoc networks the energy level of the node can be maximized.

IV ERCD PROTOCOL

In this section, we introduce our distributed clone detection protocol, namely ERCD protocol, which can achieve a high clone detection probability with little negative impact on network lifetime and limited requirement of buffer storage capacity. The ERCD protocol consists of two stages: witness selection and legitimacy verification. In witness selection, a random mapping function is employed to help each source node randomly select its witnesses. In the legitimacy verification, a verification request is sent from the source node to its witnesses, which contains the private information of the source node. If witnesses receive the verification messages, all the messages will be forwarded to the witness header for legitimacy verification, where witness headers are nodes responsible for determining whether the source node is legitimacy or not by comparing the messages collected from all witnesses. If the received messages are different from existing



record or the messages are expired, the witness header will report a lone attack to the sink to trigger a revocation procedure. Initially, the network region is virtually divided into h adjacent rings, where each ring has a sufficiently large number of sensor nodes to forward along the ring and the width of each ring is r . To simplify the description we use hop length to represent the minimal number of hops in the paper. Since we consider a densely deployed WSN, hop length of the network is the quotient of the distance from the sink to the sensor at the border of network region over the transmission range of each sensor, i.e., the distance of each hop refers to the transmission range of sensor nodes.

The ERCD protocol starts with a breadth-first search by the sink node to initiate the ring index, and all neighbouring sensors periodically exchange the relative location and ID information. After that, whenever a sensor node establishes a data transmission to others, it has to run the ERCD protocol, i.e., witness selection and legitimacy verification, to verify its legitimacy. In witness selection, a ring index is randomly selected by the mapping function as the witness ring of node a . To help relieve the traffic load in hot spot, the area around the sink cannot be selected by the mapping function. After that, node a sends its private information

to the node located in witness ring, and then the node forwards the information along the witness ring to form a ring structure. In the legitimacy verification, a verification message of the source node is forwarded to its witnesses. The ring index of node a , denoted O_a , is compared with its witness ring index O_a

w to determine the next forwarding node. If $O_a > O_w$, the message will be forwarded to any node located in ring $O_a + 1$; otherwise, the message will be forwarded to any node in ring $O_a - 1$. This step can forward the message toward the witness ring of node a . The ERCD protocol repeats above operations until a node, denoted b , located in the witness ring O_a is reached. Node b stores the private information of node a and forwards the message to any node located in ring O_a within its transmission range, denoted as c . Then, node c stores the information and forwards the message to the node d , where link (c, d) has longest projection on the extension line of the directional link from b to c . The procedure will be repeated until node b reappears in the transmission range. Therefore, the witnesses of node a have a ring structure

In the legitimacy verification, node a sends a verification message including its private information following the same path towards the witness ring as in witness selection. To enhance the probability that witnesses can

successfully receive the verification message for clone detection, the message will be broadcast when it is very close to the witness ring, namely three-ring broadcasts, i.e., the message will be broadcast in O_a

V PERFORMANCE ANALYSIS

In this section, the performance of the ERCD protocol is evaluated in terms of clone detection probability, power consumption, network lifetime, and data buffer capacity. At first, we prove that the clone detection probability of the ERCD protocol can almost surely achieve probability 1 under the scenario that witnesses are trustful in Subsection V-A. Then, we derive the expression of energy consumption and network lifetime by using ERCD protocol, and obtain the ratio of network lifetime by using ERCD protocol over RED or LSM protocol in Subsection V-B. Finally, the required data buffer of the ERCD protocol is derived in Subsection V-C.

Probability of Clone Detection

In distributed clone detection protocol with random witness selection, the clone detection probability generally refers to whether witnesses can successfully receive the verification message from the source node or not. Thus, the clone detection probability of ERCD protocol is the probability that the verification message can be successfully

transmitted from the source node to its witnesses. In ERCD protocol, the verification message is broadcast when it is near the witness ring, i.e., in the rings security. With such kind of method and assumption of trustful witnesses, we can prove that at least one of the witnesses can receive the message, i.e., the clone attack can be detected with probability one. To simplify the analysis, the transmission ranges of all sensor nodes, r , are the same.

VI CONCLUSION

In this paper, we have proposed distributed energy efficient clone detection protocol with random witness selection. Specifically, we have proposed the ERCD protocol, which includes the witness selection and legitimacy verification stages. Both of our theoretical analysis and simulation results have demonstrated that our protocol can detect the clone attack with almost probability 1, since the witnesses of each sensor node is distributed in a ring structure which makes it easy to be achieved by verification message. In addition, our protocol can achieve better network lifetime and total energy consumption with reasonable storage capacity of data buffer. This is because we take advantage of the location information by distributing the traffic load all over WSNs, such that the energy consumption and memory storage of the sensor nodes around the sink node



can be relieved and the network lifetime can be extended. In our future work, we will consider different mobility patterns under various network scenarios.

REFERENCES

- [1] Z. Zheng, A. Liu, L. X. Cai, Z. Chen, and X. Shen, "ERCD: An energy-efficient clone detection protocol in wsns," in Proc. IEEE INFOCOM, Turin, IT, Apr. 14-19 2013, pp. 2436-2444.
- [2] R. Lu, X. Li, X. Liang, X. Shen, and X. Lin, "GRS: The green, reliability, and security of emerging machine to machine communications," IEEE Communications Magazine, vol. 49, no. 4, pp. 28-35, Apr. 2011.
- [3] Christo Ananth, A. Nasrin Banu, M. Manju, S. Nilofer, S. Mageshwari, A. Peratchi Selvi, "Efficient Energy Management Routing in WSN", International Journal of Advanced Research in Management, Architecture, Technology and Engineering (IJARMATE), Volume 1, Issue 1, August 2015, pp: 16-19
- [4] Liu, J. Ren, X. Li, Z. Chen, and X. Shen, "Design principles and improvement of cost function based energy aware routing algorithms for wireless sensor networks," Computer Networks, vol. 56, no. 7, pp. 1951-1967, May. 2012.
- [5] T. Shu, M. Krunz, and S. Liu, "Secure data collection in wireless sensor networks using randomized dispersive routes," IEEE Transactions on Mobile Computing, vol. 9, no. 7, pp. 941-954, Jul. 2010.
- [6] P. Papadimitratos, J. Luo, and J. P. Hubaux, "A randomized countermeasure against parasitic adversaries in wireless sensor networks,"
- [7] IEEE Journal on Selected Areas in Communications, vol. 28, no. 7, pp. 1036-1045, Sep. 2010.
- [8] R. Lu, X. Lin, T. H. Luan, X. Liang, and X. Shen, "Pseudonym changing at social spots: An effective strategy for location privacy in VANETs," IEEE Transactions on Vehicular Technology, vol. 61, no. 1, pp. 86-96, Jan. 2012.
- [9] Z. M. Fadlullah, M. Fouda, N. Kato, X. Shen, and Y. Nozaki, "An early warning system against malicious activities for smart grid communications," IEEE Network, vol. 25, no. 5, pp. 50-55, May. 2011.
- [10] R. Lu, X. Lin, X. Liang, and X. Shen, "A dynamic privacy-preserving key management scheme for location based services in VANETs," IEEE Transactions on Intelligent Transportation Systems, vol. 13, no. 1, pp. 127-139, Jan. 2012.
- [11] M. Conti, R. D. Pietro, L. Mancini, and A. Mei, "Distributed detection of clone attacks in wireless sensor networks," IEEE Transactions on



Dependable and Secure Computing, vol. 8, no. 5, pp. 685–698, Sep.-Oct. 2011.

[12] Parno, A. Perrig, and V. Gligor, “Distributed detection of node replication attacks in sensor networks,” in Proc. IEEE Symposium on Security and Privacy, Oakland, CA, USA, May. 8-11 2005, pp. 49–63.

[13] Y. Zeng, J. Cao, S. Zhang, S. Guo, and L. Xie, “Random-walk based approach to detect clone attacks in wireless sensor networks,” IEEE Journal on Selected Areas in Communications, vol. 28, no. 28, pp. 677–691, Jun. 2010.

[14] Zhu, S. Setia, S. Jajodia, S. Roy, and L. Wang, “Localized multicast: Efficient and distributed replication detection in large-scale sensor networks,” IEEE Transactions on Mobile Computing, vol. 9, no. 7, pp. 913–926, Jul. 2010.

[15] Y. Xuan, Y. Shen, N. P. Nguyen, and M. T. Thai, “A trigger identification service for defending reactive jammers in WSN,” IEEE Transactions on Mobile Computing, vol. 11, no. 5, pp. 793–806, May. 2012.

[16] R. Lu, X. Lin, H. Zhu, X. Liang, and X. Shen, “BECAN: A bandwidth-efficient cooperative authentication scheme for filtering injected false data in wireless sensor networks,” IEEE Transactions on Parallel and Distributed Systems, vol. 23, no. 1, pp. 32–43, Jan. 2012.

[17] J. Li, J. Chen, and T. H. Lai, “Energy-efficient intrusion detection with a barrier of

probabilistic sensors,” in Proc. IEEE INFOCOM, Orlando, FL, USA, Mar. 25-30 2012, pp. 118–126